Reducing the Run-time of MCMC Programs by Multithreading on SMP Architectures

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Outline



Markov Chain Monte Carlo

- Introduction to MCMC
- Program Cycle
- Existing Parallelisation

2 Speculative Moves

- Method
- Theoretical Results
- Practical Results



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Introduction to MCMC Program Cycle Existing Parallelisation

What is Markov Chain Monte Carlo?

- MCMC is a computationally expensive iterative technique for sampling from a probability distribution.
- Basic idea:
 - Construct a Markov Chain such that its stationary distribution is equal to the distribution we wish to sample.
 - After sufficient burn-in time, sampling from the chain is equivalent to sampling from the distribution.

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Introduction to MCMC Program Cycle Existing Parallelisation

What uses Markov Chain Monte Carlo?

MCMC is widely used in

- Bayesian statistics
- Computational physics
- Computational biology
- Specific applications include:
 - Phylogenetic analysis
 - Spectral modelling of X-ray data from the Chandra X-ray satellite

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- Calculating financial econometrics
- Mapping vascular trees from retinal slides

Introduction to MCMC Program Cycle Existing Parallelisation

The MCMC Program Cycle





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Existing Parallelisation

• Execute multiple chains. Take samples from all of them.

- Embarrassingly parallel
- Does not reduce burn-in time.
- Does not help escape local optima.
- Metropolis-Coupled MCMC
 - Execute multiple chains.
 - Coarse parallelisation, machines connected by LAN.
 - Modifies algorithm to improve rate of convergence.
 - Good for escaping local optima.
 - Hard to predict benefits.

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Method Theoretical Results Practical Results

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- Each state in a Markov Chain must depend on only the preceding state.
- 2 But, typically only $\frac{1}{4}$ of iterations accept the proposed state-change.
- Consecutive rejected iterations could have been performed in parallel.
- In the second second

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- Sut, typically only $\frac{1}{4}$ of iterations accept the proposed state-change.
- Consecutive rejected iterations could have been performed in parallel.
- Sector 2 => Assume all iterations will be rejected.

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Speculative Moves Program Cycle



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Speculative Moves Program Cycle



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Theoretical Benefits of Speculative Moves

- Let:
 - *n* be the number of iterations considered concurrently.
 - *p*_r be the average state-change rejection probability.
- Each program cycle performs 1..n MCMC iterations.
- On average $\frac{1-p_{\ell}^{p}}{1-p_{r}}$ MCMC iterations are performed at each loop of the program cycle.
- If multithreading overhead negligible, time for 1 program cycle \approx time for 1 MCMC iteration.

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Theoretical Results



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Theoretical Results



Maximum benefit of speculative moves on runtime

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Practical Testing

- Circle detection algorithm used for testing
 - Fixed number of iterations
 - Autogenerated images
 - Runtime values averaged over 20 runs
- Hardware utilised:
 - AMD Athlon 64 X2 4400+ (dual-core)
 - Intel Xeon Dual-Processor
 - Intel Pentium-D (dual core)
 - Intel Core2 Quad Q6600 (2x dual-core dies)
 - 56 Itanium2 processor SGI Altix

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Comparing Practical with Theoretical (1)



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Comparing Practical with Theoretical (2)



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Preferable Architectures



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Multithreading MCMC

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This table shows the iteration time at which the overhead from multithreading balances the benefits, when $p_r = 0.75$.

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	Iteration	Iteration
	Time (µ <i>s</i>)	Rate (<i>s</i> ⁻¹)
Xeon Dual-Processor	70	14 285
Pentium-D (dual core)	55	18 181
Q6600 (using 2 threads)	75	13 333
Q6600 (using 4 threads)	25	40 000



- The speculative moves method uses increasingly available multiprocessor and multicore machines to reduce the runtime of MCMC program.
- The statistical algorithm is preserved. Speculative moves will not effect the results, only the real-time required to obtain them.
- Real-time reductions of 35% using a dual-core and 55% using quad-cores machines have been demonstrated.

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